Çağrı Çöltekin ccoltekin@sfs.uni-tuebingen.de - Lists - Stacks - Oueses Winter Semester 2023-2024 Abstract data types and data structures Arrays An array is simply a contiguous sequence of objects with the same size . An abstract data type (ADT), or abstract data structure, is an object with · Arrays are very close to how computers str well-defined operations. For example a stack supports push () and pop () data in their memory a[1] operations · Arrays can also be multi-dimensional. For An abstract data type can be implemented using different data struc example a stack can be implemented using a linked list, or an array example, matrices can be represented with 2-dimensional arrays . Sometimes the names and their usage are confusingly similar · Arrays support fast access to their elements through indexing On the downside, resizing and inserting values in arbitrary locations are expensive a = [3, 6, 8, 9, 3, 0] Arrays Liete · Main operations for list ADT are - append (and prepend)
- head (and tail) · Lists are typically implen ted using linked lists (but array-b a = [3, 6, 8, 9, 3, 0] a[0] # 3 · No built-in array data also comm head -3 6 8 9 3 0 structure in Python a[-1] . Lists are indevable For proper/faster arrays, use the numpy library a2d = [[3, 6, 8], [9, 3, 0]] a2d[0][1] # 6 head 3 6 8 9 3 0 · Python lists are array-based Stacks Oueues . A queue is a first-in-first (FIFO) out data structure A stack is a last-in-first (LIFO) out data str Two basic op Two basic operations: - enqueue - dequeu - push - pop Queues can be implemented using linked lists (or maybe arrays) . Stacks can be implemented using linked lists (or arrays) enqueue(3) Other common ADT Studying algorithms . In this course we will study a series of important algorithms, including Strings are often implemented based on character arrays
 Maps or dictionaries are similar to arrays and lists, but allow indexing with (almost) arbitrary data types Sorting
Pattern matching
Graph traversal For any algorithm we design/use, there are a number of desirable properties
Correctness an algorithm should dow that it is supposed to do
Robustness an algorithm should downer(by) handle all possible inputs it may receive
Efficiency an algorithm should be simple as possible inputs it may receive
Simplicity an algorithm should be a simple as possible - Maps are generally implemented using hashing (later in this course) Sets implement the mathematical (finite) sets: a collection unique elements without order Trees are used in n any algorithms we discuss later (we will revisit trees as data structures) \* We will briefly touch upon a few of these issues with a simple case study A simple problem: searching a sequence for a value Linear search: take 2 def linear\_search(seq, val):
answer = None linear\_search(seq, val):
for i in range(len(seq)):
 if seq[i] == val:
 return i
return None for i in range(len(seq)): if seq[i] == val: amswer = i Can we do even better? Is this a good algorithm? Can we improve it?

Recap: basic data structures

Data Structures and Algorithms for Computational Linguistics III

ISCL-BA-07

Linear search: take 3 Binary search def linear\_search(seq, val): n = len(seq) - i idef binary\_search(seq, val):
 left, right = 0, len(seq)
 while left <= right:
 mid = (left +right) // 2
 if seq[mid] == val:
 return mid
 f seq[mid] > val: n = len(seq) - 1 last = seq[n] seq[n] = val i = 0 while seq[i] != val: i += 1 Is this better? · Any disadvantages? the sequence is sorted. i += 1
seq[n] = last
if i < n or seq[n] == val:
 return i
else:
 return None</pre> . Can we do even better? right = mid - 1 left = mid + 1 return None Binary search A note on recursion · Some problems are much easier to solve recursively. def binary\_search\_recursive(seq, val, left=None, right=None): Recursion is also a mathematical concept, properties of recu are often easier to prove if left is Nor left = 0 if right is None Reminder: right = len(seq)
if left > right: - You have to define one or more base cases (e.g., if left > right for binary search)

- Each recursive step should approach the base case (e.g., should run on a smaller portion of the data) return Nome

nid = (left + right) // 2

if seq[nid] = val:
 return mid

if seq[nid] val:
 return binary\_search\_recursive(seq, val, left, mid - 1) \* We will see quite a few recursive algorithms, it is time for getting used to if you are not else: return binary\_search\_recursive(seq, val, mid + 1, right) Exercise: write a recursive function for linear search. Summary An interesting historical anecdote How difficult is the 'binary search'? This lecture is a review of some basic data structure and algorith 1946 suggested in a lecture by John Mauchly We will assume you know these concepts, please revise your earlier knowledge if needed 1960 first fix to the original version (by Derrick Henry Lehmer) 1962 another fix/improvement was published (by Hermann Bottenbruch) Next: 2006 a bug was discovered in Java's binary search implementation Analysis of algorithms (Reading: textbook (Goodrich, Tamassia, and Goldwasser, 2013) chapter 3) . A few common patterns for designing (efficient) algorithms Although the basic idea of binary search is comparatively straightforward the details can be surprisingly tricky. -Knuth (1973) Acknowledgments, credits, references . Some of the slides are based on the previous year's course by Corina Dima. . The historical information on binary search development on Slide 1 is from Wikipedia Goodrich, Michael T., Roberto Tamassia, and Michael H. Goldwasser (2013). Data Structures and Algorithms in Python. John Wiley & Sons, Incorporated. 158 2.011184/6/34.
Knuth, Donald E. (1973). The Art of Computer Programming: Sorting and Searching. Vol. 3. Pearson Education. sinc. 9780321633785.