Bottom-up Chart Parsing: the CKY algorithm Data Structures and Algorithms for Computational Linguistics III (ISCL-BA-07)

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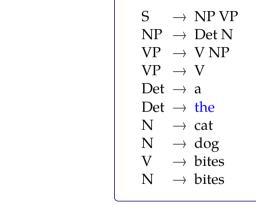
Parsing so far

- Parsing is the task of automatic syntactic analysis
- For most practical purposes, context-free grammars are the most useful formalism for parsing
- We can formulate parsing as
 - Top-down: begin with the start symbol, try to *produce* the input string to be parsed
 - Bottom up: begin with the input, and try to *reduce* it to the start symbol
- Both strategies can be cast as search with backtracking
- Backtracking parsers are inefficient: they recompute sub-trees multiple times

S \rightarrow NP VP $NP \rightarrow Det N$ $VP \rightarrow VNP$ $VP \rightarrow V$ Det \rightarrow a Det \rightarrow the $N \rightarrow cat$ $N \rightarrow dog$ $V \rightarrow bites$ $N \rightarrow bites$

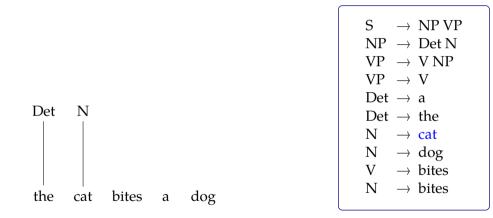
the cat bites a dog

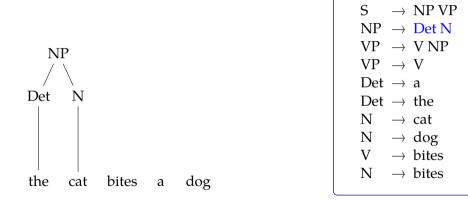
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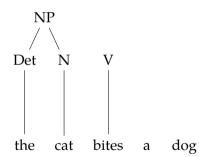


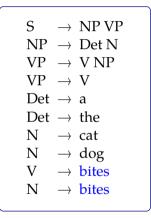
the cat bites a dog

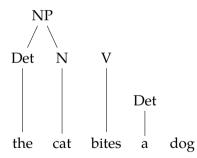
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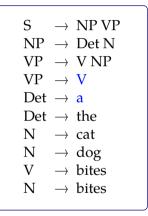


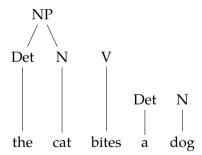






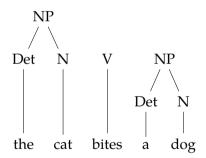


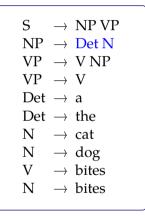


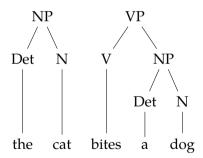


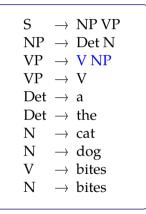
$$\begin{array}{rcl} S & \rightarrow & NP \ VP \\ NP & \rightarrow & Det \ N \\ VP & \rightarrow & V \ NP \\ VP & \rightarrow & V \\ Det & \rightarrow & a \\ Det & \rightarrow & the \\ N & \rightarrow & cat \\ N & \rightarrow & dog \\ V & \rightarrow & bites \\ N & \rightarrow & bites \end{array}$$

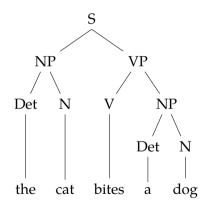
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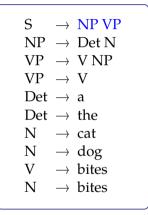












Introduction CNF CKY

Dealing with ambiguity

S \rightarrow NP VP $NP \rightarrow Prn N$ $NP \rightarrow Prn$ $VP \rightarrow V NP$ $VP \rightarrow V$ $VP \rightarrow VS$ $N \rightarrow duck$ V \rightarrow duck V \rightarrow saw $\Pr n \rightarrow I$ \leftarrow $Prn \rightarrow she$ $Prn \rightarrow her$

I saw her duck

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Introduction CNF CKY

Dealing with ambiguity

Prn

S \rightarrow NP VP $NP \rightarrow Prn N$ $NP \rightarrow Prn$ \leftarrow $VP \rightarrow V NP$ $VP \rightarrow V$ $VP \rightarrow VS$ $N \rightarrow duck$ $V \rightarrow duck$ V \rightarrow saw $\Pr n \rightarrow I$ $Prn \rightarrow she$ $Prn \rightarrow her$

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saw

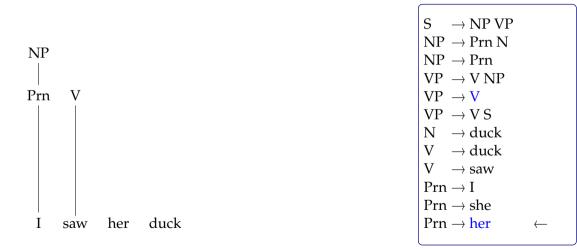
her

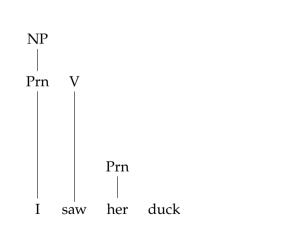
duck

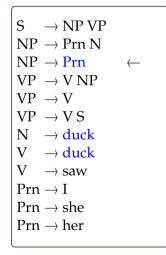
NP | Prn

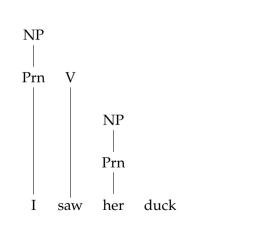
S \rightarrow NP VP $NP \rightarrow Prn N$ $NP \rightarrow Prn$ $VP \rightarrow V NP$ $VP \rightarrow V$ $VP \rightarrow VS$ $N \rightarrow duck$ V \rightarrow duck V \rightarrow saw \leftarrow $\Pr n \rightarrow I$ $Prn \rightarrow she$ $Prn \rightarrow her$

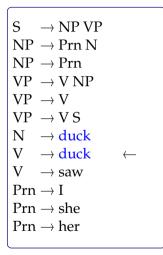
saw her duck

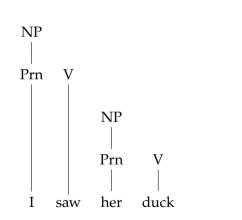


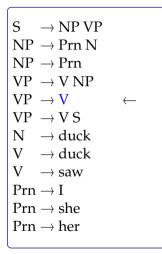


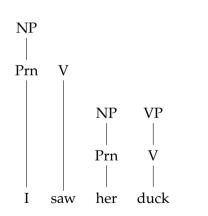




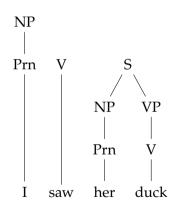




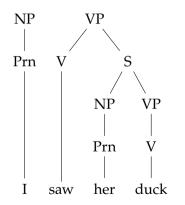


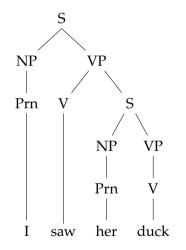


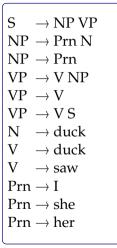
S	\rightarrow	NP VP	\leftarrow	
NP	\rightarrow	Prn N		
NP	\rightarrow	Prn		
VP	\rightarrow	V NP		
VP	\rightarrow	V		
VP	\rightarrow	VS		
Ν	\rightarrow	duck		
V	\rightarrow	duck		
V	\rightarrow	saw		
Prn	\rightarrow	Ι		
Prn	\rightarrow	she		
Prn	\rightarrow	her		

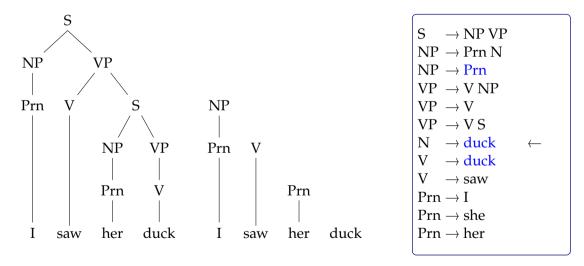


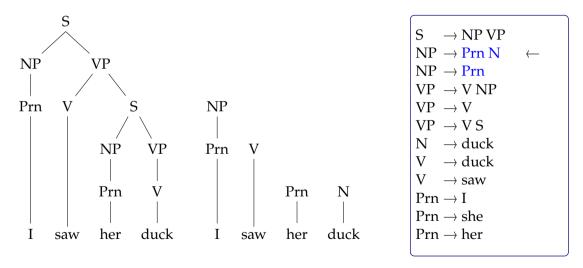
S	\rightarrow	NP VP	
NP	\rightarrow	Prn N	
NP	\rightarrow	Prn	
VP	\rightarrow	V NP	
VP	\rightarrow	V	
VP	\rightarrow	VS	\leftarrow
Ν	\rightarrow	duck	
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Prn	\rightarrow	her	

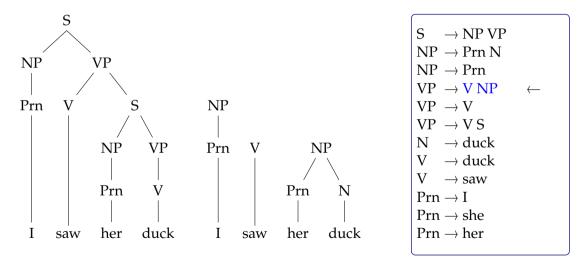


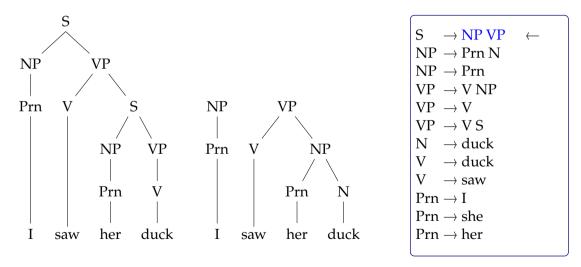


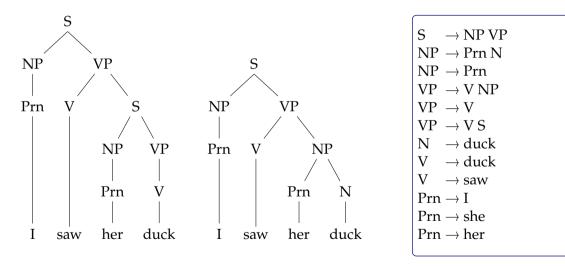








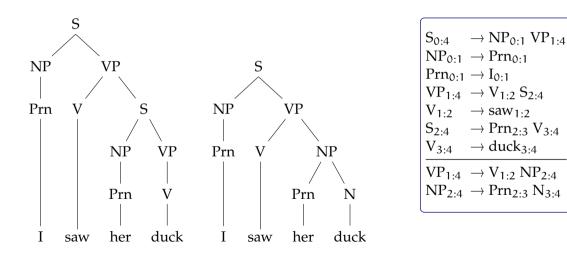




Introduction CNF CKY

How to represent multiple parses

parse forest grammar



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CKY algorithm

- The CKY (Cocke–Kasami–Younger) parsing algorithm is a dynamic programming algorithm
- It processes the input *bottom up*, and saves the intermediate results on a *chart*
- Time complexity for *recognition* is $O(n^3)$
- Space complexity is $O(\mathfrak{n}^2)$
- It requires the CFG to be in *Chomsky normal form* (CNF) (can somewhat be relaxed, but not common)

Chomsky normal form (CNF)

- A CFG is in CNF, if the rewrite rules are in one of the following forms
 - $\ A \ \rightarrow B \ C$
 - $\ A \ \rightarrow a$

where A, B, C are non-terminals and a is a terminal

- Any CFG can be converted to CNF
- Resulting grammar is *weakly equivalent* to the original grammar:
 - it generates/accepts the same language
 - but the derivations are different

S S NP VP VP	\rightarrow \rightarrow	NP VP Aux NP VP the N V NP V
Ν	\rightarrow	cat
Ν	\rightarrow	dog
V	\rightarrow	bites
Ν	\rightarrow	bites

 $S \rightarrow NP VP$ $S \rightarrow Aux NP VP$ $NP \rightarrow the N$ $VP \rightarrow VNP$ $VP \rightarrow V$ $N \rightarrow cat$ $N \rightarrow dog$ V \rightarrow bites $N \rightarrow bites$

•
$$S \rightarrow Aux NP VP$$

 $S \rightarrow Aux NP VP \Rightarrow S \rightarrow Aux X$
 $X \rightarrow NP VP$

 $S \rightarrow NP VP$ $S \rightarrow Aux NP VP$ $NP \rightarrow the N$ $VP \rightarrow VNP$ $VP \rightarrow V$ $N \rightarrow cat$ $N \rightarrow dog$ V \rightarrow bites $N \rightarrow bites$

•
$$S \rightarrow Aux NP VP$$

 $S \rightarrow Aux NP VP \Rightarrow S \rightarrow Aux X$
 $X \rightarrow NP VP$
• $NP \rightarrow the N$
 $NP \rightarrow the N \Rightarrow NP \rightarrow X N$
 $X \rightarrow the$

 $S \rightarrow NP VP$ $S \rightarrow Aux NP VP$ $NP \rightarrow the N$ $VP \rightarrow VNP$ $VP \rightarrow V$ $N \rightarrow cat$ $N \rightarrow dog$ V \rightarrow bites $N \rightarrow bites$

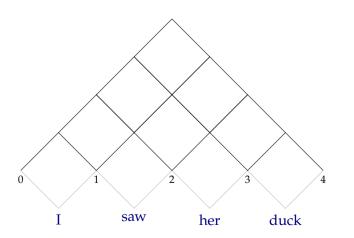
•
$$S \rightarrow Aux NP VP$$

 $S \rightarrow Aux NP VP \Rightarrow S \rightarrow Aux X$
 $X \rightarrow NP VP$
• $NP \rightarrow the N$
 $NP \rightarrow the N \Rightarrow NP \rightarrow X N$
 $X \rightarrow the$
• $VP \rightarrow V$
 $VP \rightarrow V \Rightarrow VP \rightarrow bites$

Converting to CNF

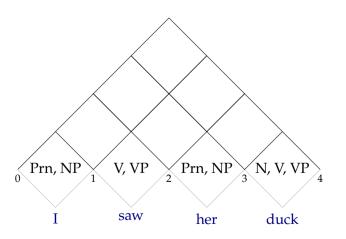
- 1. Eliminate the ε rules: if $A \ \rightarrow \varepsilon$ is in the grammar
 - replace any rule $B \ \rightarrow \alpha \ A \ \beta$ with two rules
 - $\begin{array}{l} B \ \rightarrow \alpha \ \beta \\ B \ \rightarrow \alpha \ A' \ \beta \end{array}$
 - add $A' \to \alpha$ for all α (except $\varepsilon)$ whose LHS is A
 - repeat the process for newly created ϵ rules
 - remove the rules with ε on the RHS (except S $\rightarrow \varepsilon)$
- 2. Eliminate unit rules: for a rule A \rightarrow B
 - Replace the rule with A $\rightarrow \alpha_1 \mid ... \mid \alpha_n$, where $\alpha_1, ..., \alpha_n$ are all RHS or rule B
 - Remove the rule A \rightarrow B
 - Repeat the process until no unit rules remain
- 3. Binarize all the non-binary rules with non-terminal on the RHS: for a rule $A \ \rightarrow \ X_1 \ X_2 \ ... \ X_n :$
 - Replace the rule with A $\ \rightarrow$ A1 X3...Xn, and add A1 $\ \rightarrow$ X1 X2
 - Repeat the process until all new rules are binary

an ambiguous example



$$\begin{array}{ll} S & \rightarrow NP \, VP \\ NP & \rightarrow Prn \, N \\ VP & \rightarrow V \, NP \\ VP & \rightarrow V \, S \\ N & \rightarrow duck \\ VP & \rightarrow duck \, | \, saw \\ V & \rightarrow duck \, | \, saw \\ Prn & \rightarrow I \, | \, she \, | \, her \\ NP & \rightarrow I \, | \, she \, | \, her \end{array}$$

an ambiguous example



an ambiguous example

NP VP S \rightarrow ? \bigvee_{3} N, V, VP Prn, NP V, VP Prn, NP 0 2 saw duck her

 $\begin{array}{ll} S & \rightarrow NP \ VP \\ NP & \rightarrow Prn \ N \\ VP & \rightarrow V \ NP \\ VP & \rightarrow V \ S \\ N & \rightarrow duck \end{array}$

$$VP \rightarrow duck \mid saw$$

$$V \rightarrow duck \mid saw$$

$$\Pr{} \to I \mid she \mid her$$

 $NP \, \rightarrow I \, | \, she \, | \, her$

Introduction CNF CKY

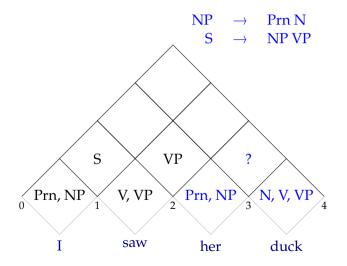
CKY demonstration

an ambiguous example

V NP VP \rightarrow S ? Prn, NP V, VP Prn, NP N, V, VP 0 2 saw duck her

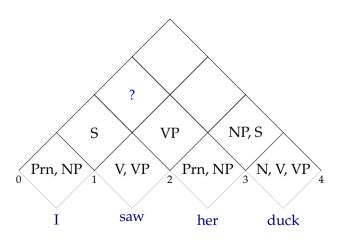
- $\begin{array}{ll} S & \rightarrow NP \ VP \\ NP & \rightarrow Prn \ N \\ VP & \rightarrow V \ NP \\ VP & \rightarrow V \ S \\ N & \rightarrow duck \\ VP & \rightarrow duck \ | \ saw \\ V & \rightarrow duck \ | \ saw \\ Prn & \rightarrow I \ | \ she \ | \ her \end{array}$
- $NP \rightarrow I | she | her$

an ambiguous example



 $\begin{array}{ll} S & \rightarrow NP \ VP \\ NP & \rightarrow Prn \ N \\ VP & \rightarrow V \ NP \\ VP & \rightarrow V \ S \\ N & \rightarrow duck \\ VP & \rightarrow duck \ | \ saw \\ VP & \rightarrow duck \ | \ saw \\ Prn & \rightarrow I \ | \ she \ | \ her \\ NP & \rightarrow I \ | \ she \ | \ her \end{array}$

an ambiguous example



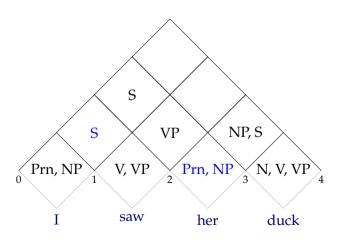
an ambiguous example

NP VP S S NP, S S VP Prn, NP V, VP Prn, NP N, V, VP 0 2 2 saw duck her

 $\begin{array}{ll} S & \rightarrow NP \ VP \\ NP & \rightarrow Prn \ N \end{array}$

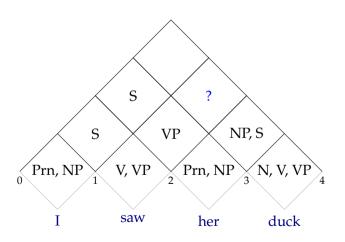
- $VP \ \rightarrow V \ NP$
- $VP \ \rightarrow V \ S$
- $N \quad \rightarrow duck$
- $VP \ \rightarrow duck \,|\, saw$
- $V \rightarrow duck \mid saw$
- $Prn \to I \ | \ she \ | \ her$
- $NP \, \rightarrow I \, | \, she \, | \, her$

an ambiguous example



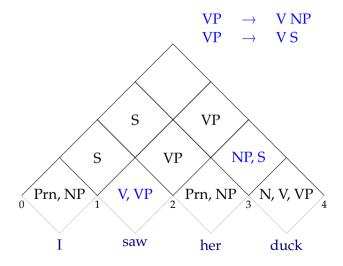
 $\begin{array}{ll} S & \rightarrow NP \ VP \\ NP & \rightarrow Prn \ N \\ VP & \rightarrow V \ NP \\ VP & \rightarrow V \ S \\ N & \rightarrow duck \\ VP & \rightarrow duck \ | \ saw \\ V & \rightarrow duck \ | \ saw \\ Prn & \rightarrow I \ | \ she \ | \ her \\ NP & \rightarrow I \ | \ she \ | \ her \end{array}$

an ambiguous example



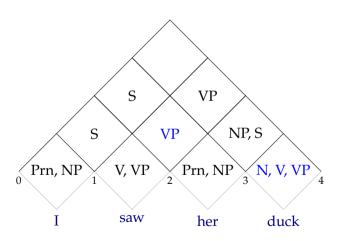
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an ambiguous example



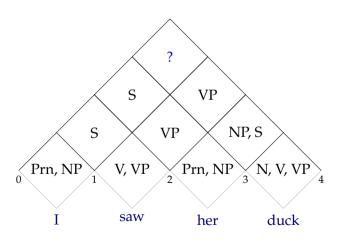
S \rightarrow NP VP $NP \rightarrow Prn N$ $VP \rightarrow V NP$ $VP \rightarrow VS$ $N \rightarrow duck$ VP \rightarrow duck | saw \rightarrow duck | saw V $Prn \rightarrow I \mid she \mid her$ $NP \rightarrow I \mid she \mid her$

an ambiguous example



S \rightarrow NP VP $NP \rightarrow Prn N$ $VP \rightarrow V NP$ $VP \rightarrow VS$ $N \rightarrow duck$ VP \rightarrow duck | saw \rightarrow duck | saw V $Prn \rightarrow I \mid she \mid her$ $NP \rightarrow I \mid she \mid her$

an ambiguous example



 $\begin{array}{ll} S & \rightarrow NP \ VP \\ NP & \rightarrow Prn \ N \\ VP & \rightarrow V \ NP \\ VP & \rightarrow V \ S \\ N & \rightarrow duck \\ VP & \rightarrow duck \ | \ saw \\ V & \rightarrow duck \ | \ saw \\ Prn & \rightarrow I \ | \ she \ | \ her \\ NP & \rightarrow I \ | \ she \ | \ her \end{array}$

an ambiguous example

NP VP S \rightarrow S S VP NP, S S VP Prn, NP V, VP Prn, NP (N, V, VP $\hat{\mathbf{x}}$ 0 2 saw duck her

 $\begin{array}{ll} S & \rightarrow NP \ VP \\ NP & \rightarrow Prn \ N \\ VP & \rightarrow V \ NP \\ VP & \rightarrow V \ S \end{array}$

$$N \rightarrow duck$$

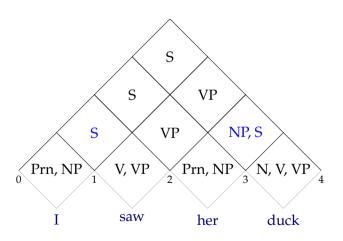
$$VP \rightarrow duck \mid saw$$

$$V \rightarrow duck \mid saw$$

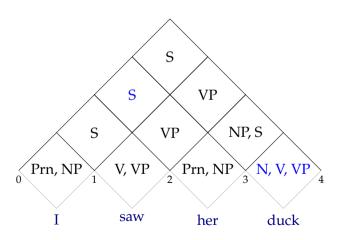
$$\Prn \to I \mid she \mid her$$

 $NP \ \rightarrow I \ | \ she \ | \ her$

an ambiguous example



an ambiguous example



 $\begin{array}{ll} S & \rightarrow NP \ VP \\ NP & \rightarrow Prn \ N \\ VP & \rightarrow V \ NP \\ VP & \rightarrow V \ S \\ N & \rightarrow duck \\ VP & \rightarrow duck \ | \ saw \\ V & \rightarrow duck \ | \ saw \\ Prn & \rightarrow I \ | \ she \ | \ her \\ NP & \rightarrow I \ | \ she \ | \ her \end{array}$

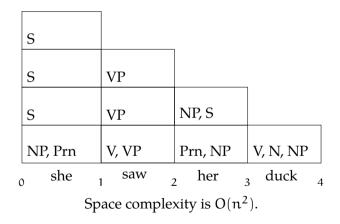
CKY demonstration: the chart

our chart is a 2D array

NP, Prn	S	S	S
	V, VP	VP	VP
		Prn	NP, S
			V, N, NP
o she	1 saw	₂ her	₃ duck ₄
Space complexity is $O(n^2)$.			

CKY demonstration: the chart

our chart is a 2D array - this is more convenient for programming

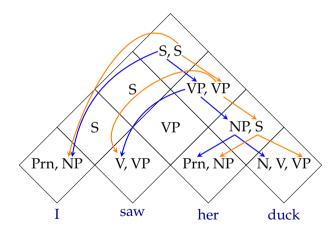


Parsing vs. recognition

- We went through a recognition example
- Note that the algorithm is not directional: it takes the complete input
- Recognition accepts or rejects a sentence based on a grammar
- For parsing, we want to know the derivations that yielded a correct parse
- To recover parse trees, we
 - follow the same procedure as recognition
 - add back links to keep track of the derivations

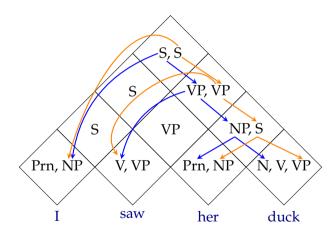
Introduction CNF CKY

Chart parsing example (CKY parsing)



Introduction CNF CKY

Chart parsing example (CKY parsing)



The chart stores a *parse forest* efficiently.

Summary

- + CKY avoids re-computing the analyses by storing the earlier analyses (of sub-spans) in a table
- It still computes lower level constituents that are not allowed by the grammar
- CKY requires the grammar to be in CNF
- CKY has $O(n^3)$ recognition complexity
- For parsing we need to keep track of backlinks
- CKY can efficiently store all possible parses in a chart
- Enumerating all possible parses have exponential complexity (worst case)
- Suggested reading: Jurafsky and Martin (2009, draft 3rd ed, section 13.2)

Summary

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Next:

- Top-down chart parsing: Earley algorithm
- Suggested reading:
 - Jurafsky and Martin (2009, section 13.2.4)
 - Grune and Jacobs (2007, section 7.2)

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Acknowledgments, references, additional reading material



Grune, Dick and Ceriel J.H. Jacobs (2007). Parsing Techniques: A Practical Guide. second. Monographs in Computer Science. The first edition is available at http://dickgrune.com/Books/PTAPG_1st_Edition/BookBody.pdf. Springer New York. ISBN: 9780387689548.

Jurafsky, Daniel and James H. Martin (2009). Speech and Language Processing: An Introduction to Natural Language Processing, Computational Linguistics, and Speech Recognition. second edition. Pearson Prentice Hall. ISBN: 978-0-13-504196-3.