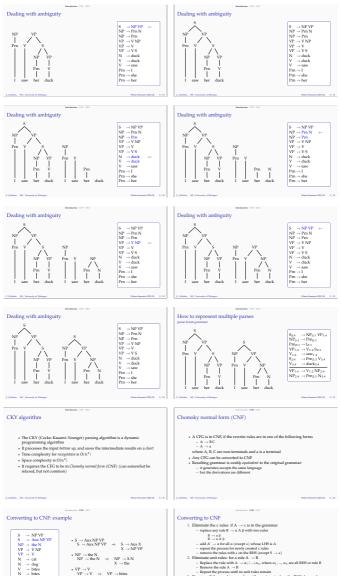
Parsing so far Bottom-up Chart Parsing: the CKY algorithm Data Structures and Algorithms for Computa (ISCL-BA-07) * Parsing is the task of automatic syntactic analysis For most practical purposes, context-free grammars are the most useful formalism for parsing
 We can formulate parsing as Çağrı Çöltekin ccoltekin@sfs.uni-tuebingen.de e can occurrence passing as

- Top-down: begin with the start symbol, try to produce the input string to be
parsed

- Bottom up: begin with the input, and try to refuce it to the start symbol · Both strategies can be cast as search with backtracking Winter Semester 2023/24 Backtracking parsers are inefficient: they recompute sub-tro Bottom-up parsing as search Dealing with ambiguity $\to NP\,VF$ $S \rightarrow NP VF$ $NP \rightarrow Det N$ $VP \rightarrow V NP$ $VP \rightarrow V$ $Det \rightarrow a$ $Det \rightarrow the$ $N \rightarrow cat$ $N \rightarrow dog$ $V \rightarrow bittes$ $N \rightarrow bittes$ $NP \rightarrow Pm N$ $NP \rightarrow Pm$ $VP \rightarrow V NP$ $VP \rightarrow V$ $VP \rightarrow VS$ N → duck V → duck V → saw $Prn \rightarrow I$ Prn → she I saw her duck Dealing with ambiguity Dealing with ambiguity → NP VP \rightarrow NP VF $S \rightarrow NP VI$ $NP \rightarrow Pm N$ $NP \rightarrow Pm$ $VP \rightarrow V NP$ $VP \rightarrow V$ $VP \rightarrow VS$ $N \rightarrow duck$ $V \rightarrow duck$ $V \rightarrow suv$ $V \rightarrow suv$ $S \rightarrow NP VF$ $NP \rightarrow Pm N$ $NP \rightarrow Pm$ $VP \rightarrow V NP$ $VP \rightarrow V S$ $N \rightarrow duck$ $V \rightarrow saw$ $Pm \rightarrow I$ $Pm \rightarrow she$ $Pm \rightarrow her$ I $Prn \rightarrow I$ $Prn \rightarrow sh$ Prn → her $Prn \rightarrow her$ Dealing with ambiguity Dealing with ambiguity → NP VP \rightarrow NP VF $NP \rightarrow Pm N$ $NP \rightarrow Pm N$ $NP \rightarrow Prn \ N$ $NP \rightarrow Prn$ $VP \rightarrow V \ NP$ $VP \rightarrow V$ $VP \rightarrow V \ S$ $V \rightarrow dtack$ $V \rightarrow saw$ $NP \rightarrow Pm N$ $NP \rightarrow Pm$ $VP \rightarrow V NP$ $VP \rightarrow V$ $VP \rightarrow V S$ $N \rightarrow duck$ $V \rightarrow duck$ $V \rightarrow saw$ Pri Prn → I Prn → she $Pm \to I$ Prn → she duck Prn → her Prn → ber Dealing with ambiguity Dealing with ambiguity \rightarrow NP VP \rightarrow NP VF NP → Prn N NP → Pm N $NP \rightarrow Pm N$ $NP \rightarrow Pm$ $VP \rightarrow V NP$ $VP \rightarrow V$ $VP \rightarrow VS$ $N \rightarrow duck$ $V \rightarrow duck$ $NP \rightarrow Pm \ N$ $NP \rightarrow Pm$ $VP \rightarrow V \ NP$ $VP \rightarrow V$ $VP \rightarrow V \ S$ $V \rightarrow duck$ $V \rightarrow saw$ V → saw Prn → I Prn → she Pm → I Pm → she Prn → her Prn → her Dealing with ambiguity Dealing with ambiguity $S \rightarrow NP VI$ $NP \rightarrow Prn N$ $S \rightarrow NP VI$ $NP \rightarrow Pm N$ $NP \rightarrow Pm N$ $NP \rightarrow Pm$ $VP \rightarrow V NP$ $VP \rightarrow V$ $VP \rightarrow V S$ $N \rightarrow duck$ $V \rightarrow duck$ $V \rightarrow saw$ $Pm \rightarrow I$ $NP \rightarrow Prn$ $NP \rightarrow Prn$ $VP \rightarrow V$ NE $VP \rightarrow V$ $VP \rightarrow V$ N → duck V → duck V → saw $\operatorname{Prn} \to \operatorname{I}$ $\operatorname{Prn} \to \operatorname{she}$ $Pm \rightarrow she$ Prn → he Prn → her



Remove the rule A → B
 Repeat the percoss until no unit rules remain
 Binarize all the non-binary rules with non-terminal on the RHS: for a rule A → X₁ X₂ x x₆;
 Replace the rule with A → A₁ X₃ x₆, and add A₁ → X₁ X₂.
 Repeat the process until all near rules are bine rules.

